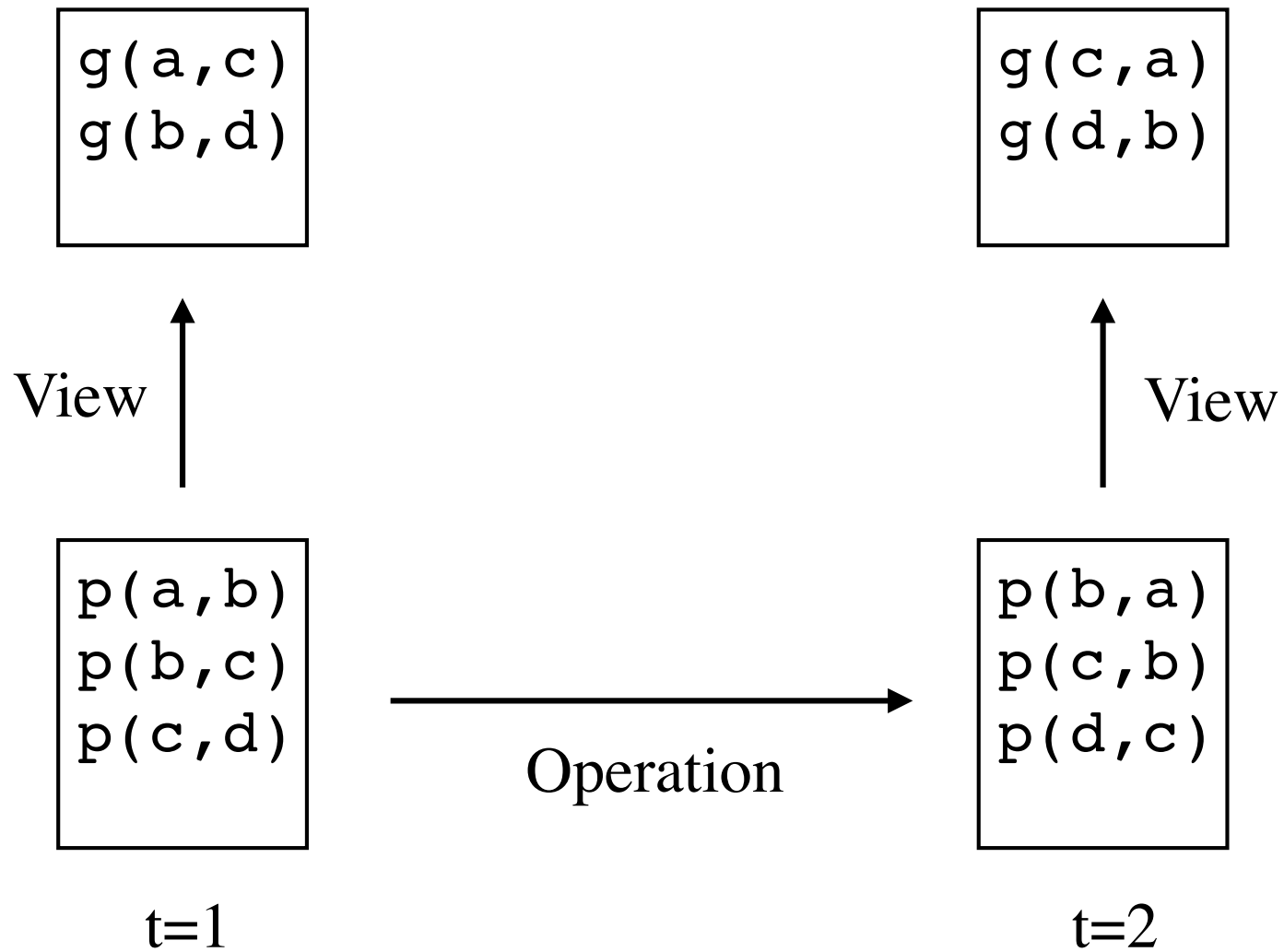


Logic Programming

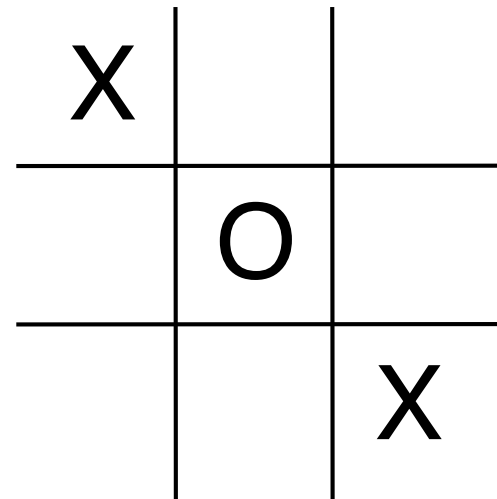
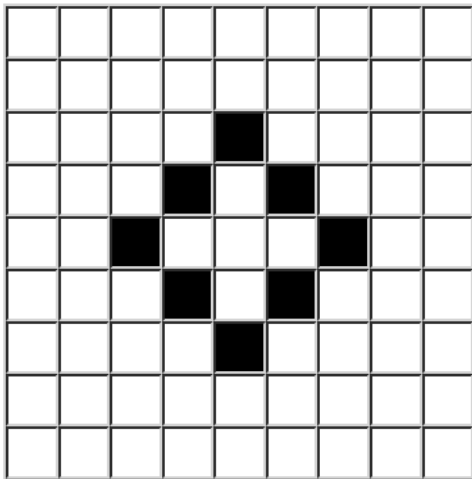
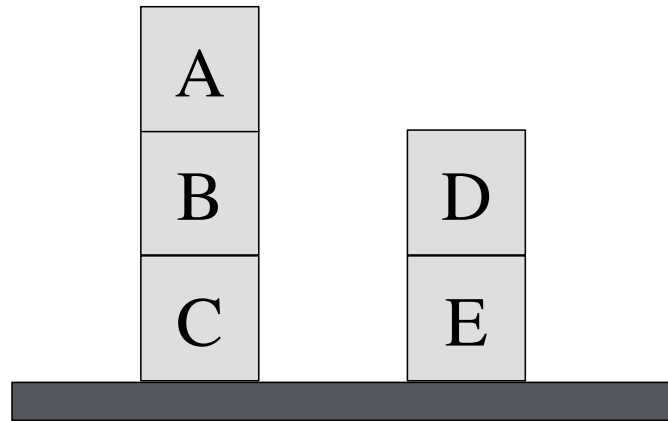
Model Management

Michael Genesereth
Computer Science Department
Stanford University

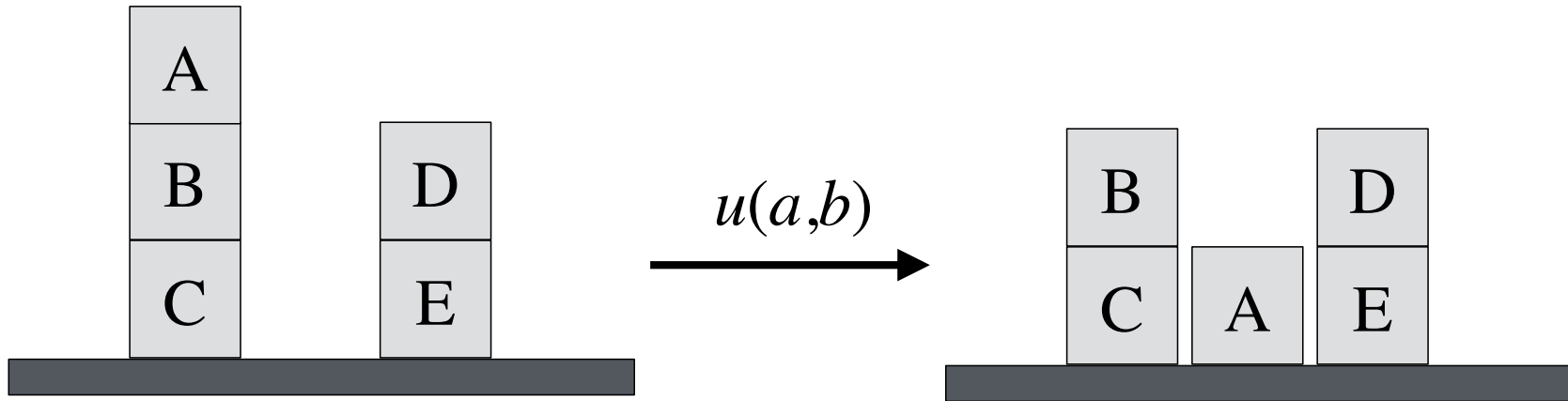
Operations



Dynamic Worlds



Changing Worlds

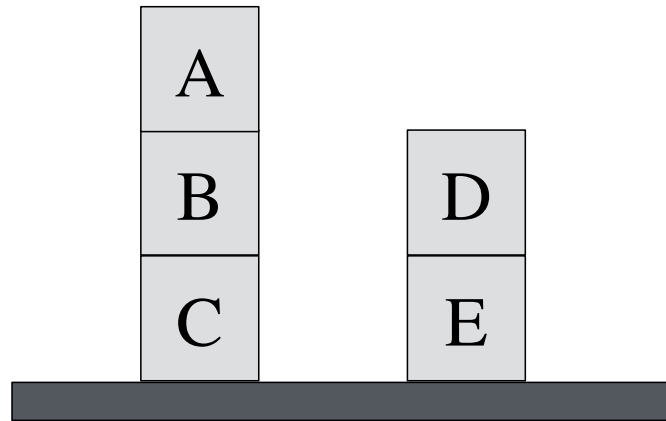


$\text{on}(a, b)$
 $\text{on}(b, c)$
 $\text{on}(d, e)$

$u(a, b)$

$\text{on}(b, c)$
 $\text{on}(d, e)$

Changing Models



```
on(a,b)
on(b,c)
on(d,e)
```

augment
→

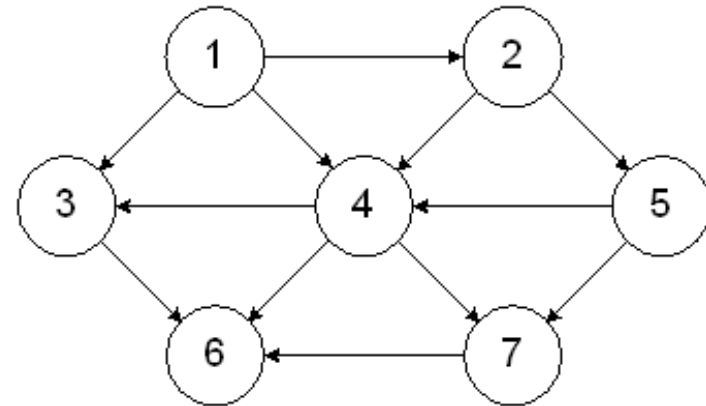
```
on(a,b)
on(b,c)
on(d,e)
clear(a)
clear(d)
table(c)
table(e)
```

Programme

Tables versus Triples

First Name	Last Name	Address	City	Age
Mickey	Mouse	123 Fantasy Way	Anaheim	73
Bat	Man	321 Cavern Ave	Gotham	54
Wonder	Woman	987 Truth Way	Paradise	39
Donald	Duck	555 Quack Street	Mallard	65
Bugs	Bunny	567 Carrot Street	Rascal	58
Wiley	Coyote	999 Acme Way	Canyon	61
Cat	Woman	234 Purrfect Street	Hairball	32
Tweety	Bird	543	Itotitaw	28

Graph Editing



View Materialization

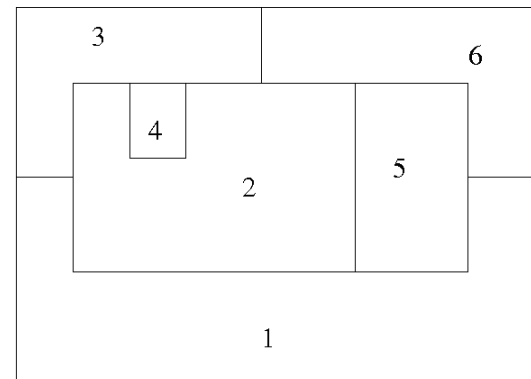
Order table					
Partition key	Row key	Order date	Shipping address	Total invoice	Order status
001 (Customer ID)	1 (Order ID)	11082013	One Microsoft way Redmond, WA 98052	\$400	In process
005	2	11082013	One Microsoft way Redmond, WA 98052	\$200	Shipped

Orderitem table					
Partition key	Row key	Product	Unit Price	Amount	Total
1 (Order ID)	001_1 (Orderitem ID)	XX	\$100	2	\$200
1	001_2	YY	\$40	5	\$200
2	002_1	ZZ	\$200	1	\$200

Customer table					
Partition key	Row key	Billing Information	Shipping address	Gender	Age
US East (region)	001 (Customer ID)	****0001	One Microsoft way Redmond, WA 98052	Female	30
US East	002	****2006	One Microsoft way Redmond, WA 98052	Male	40

Materialized View					
Partition key	Row key	Product Name	Total sold	Number of customers	
Electronics (Product category)	001 (Product ID)	XX	\$30,000	500	
Electronics	002	YY	\$100,000	400	

Constraint Propagation



Tables versus Triples

University

Students:

aaron
belinda
calvin
george

Departments:

architecture
computers
english
physics

Faculty:

alan
cathy
donna
frank

Years:

freshman
sophomore
junior
senior

Tables

Students:

aaron
belinda
calvin
george

Departments:

architecture
computers
english
physics

Faculty:

alan
cathy
donna
frank

Years:

freshman
sophomore
junior
senior

Predicate:

student (Student, Department, Advisor, Year)

Dataset:

student (aaron, architecture, alan, freshman)
student (belinda, computers, cathy, sophomore)
student (calvin, english, donna, junior)
student (george, physics, frank, senior)

Triples

```
student.major(aaron, architecture)  
student.advisor(aaron, alan)  
student.year(aaron, freshman)
```

```
student.major(belinda, computers)  
student.advisor(belinda, cathy)  
student.year(belinda, sophomore)
```

```
student.major(calvin, physics)  
student.advisor(calvin, donna)  
student.year(calvin, senior)
```

```
student.major(george, physics)  
student.advisor(george, frank)  
student.year(george, senior)
```

Convert to Triples

```
convert_to_triples ::  
  student(S,D,A,Y) ==>  
    student.major(S,D) &  
    student.advisor(S,A) &  
    student.year(S,Y)
```

```
convert_to_triples ::  
  student(S,D,A,Y) ==> ~student(S,D,A,Y)
```

- or -

```
convert_to_triples ::  
  student(S,D,A,Y) ==>  
    ~student(S,D,A,Y) &  
    student.major(S,D) &  
    student.advisor(S,A) &  
    student.year(S,Y)
```

Convert to Tables

```
convert_to_tables ::  
  student.major(S,D) &  
  student.advisor(S,A) &  
  student.year(S,Y) ==>  
  student(S,D,A,Y)
```

```
convert_to_tables ::  
  student.major(S,D) ==> ~student.major(S,D)
```

```
convert_to_tables ::  
  student.advisor(S,D) ==>  
  ~student.advisor(S,D)
```

```
convert_to_tables ::  
  student.year(S,D) ==> ~student.year(S,D)
```

Lambda

Save Revert Sort

```
student(aaron,architecture,alan,freshman)
student(belinda,computers,cathy,sophomore)
student(calvin,english,donna,junior)
student(george,physics,frank,senior)
```

Library

Save Revert

```
convert_to_triples ::
  student(S,D,A,Y) ==>
    student.major(S,D) &
    student.advisor(S,A) &
    student.year(S,Y)

convert_to_triples ::
  student(S,D,A,Y) ==> ~student(S,D,A,Y)

convert_to_tables ::
  student.major(S,D) &
  student.advisor(S,A) &
  student.year(S,Y) ==>
  student(S,D,A,Y)

convert_to_tables ::
  student.major(S,D) ==> ~student.major(S,D)

convert_to_tables ::
  student.advisor(S,D) ==>
  ~student.advisor(S,D)

convert_to_tables ::
  student.year(S,D) ==> ~student.year(S,D)
```

Execute

Action convert_to_tables

Expand Expand on update
Execute Run on clock tick

Lambda

Save Revert Sort

```
student(aaron,architecture,alan,freshman)
student(belinda,computers,cathy,sophomore)
student(calvin,english,donna,junior)
student(george,physics,frank,senior)
```

Library

Save Revert

```
convert_to_triples ::
  student(S,D,A,Y) ==>
    student.major(S,D) &
    student.advisor(S,A) &
    student.year(S,Y)

convert_to_triples ::
  student(S,D,A,Y) ==> ~student(S,D,A,Y)

convert_to_tables ::
  student.major(S,D) &
  student.advisor(S,A) &
  student.year(S,Y) ==>
    student(S,D,A,Y)

convert_to_tables ::
  student.major(S,D) ==> ~student.major(S,D)

convert_to_tables ::
  student.advisor(S,D) ==>
    ~student.advisor(S,D)

convert_to_tables ::
  student.year(S,D) ==> ~student.year(S,D)
```

Execute

Action convert_to_tables

 Expand Expand on update
 Execute Run on clock tick

Timer

Step 41

Start Pause Reset

Transform

Condition evaluate(remainder(timer(),2),0)

Conclusion convert_to_triples

 Expand Expand on update
 Execute Run on clock tick

Transform

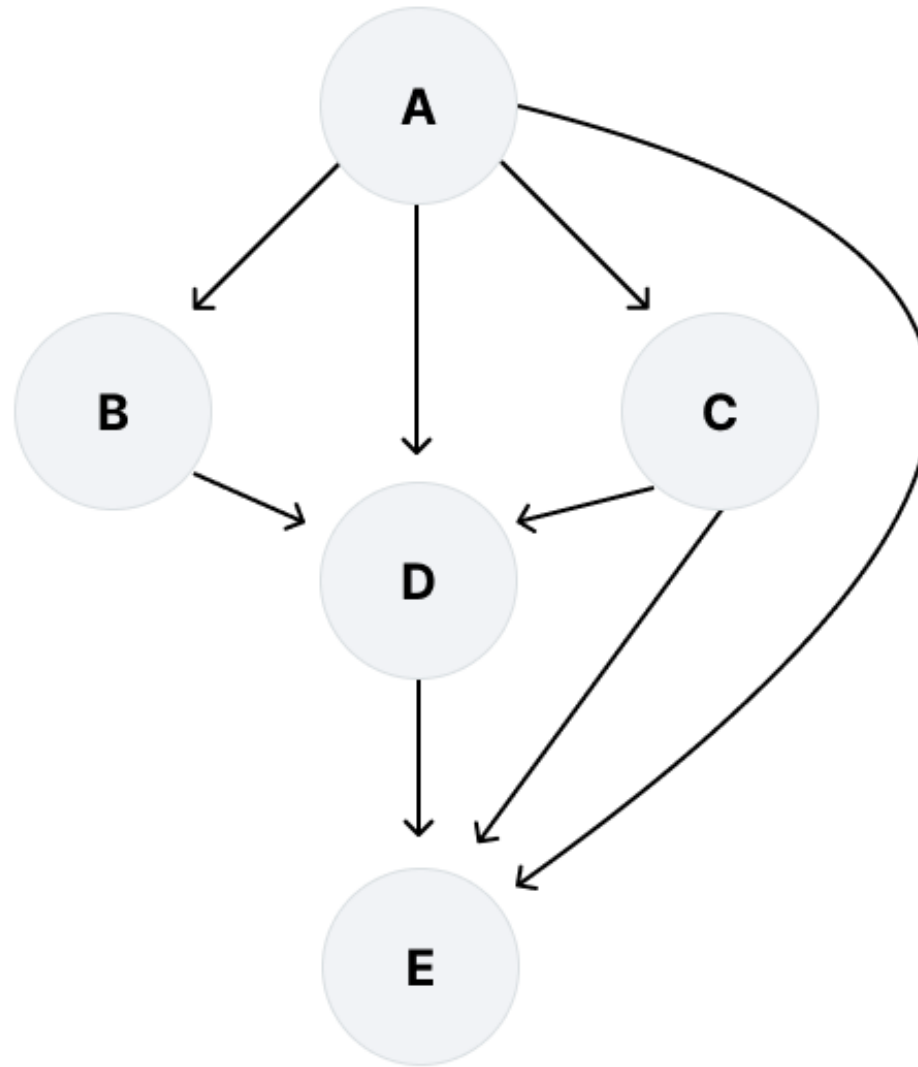
Condition evaluate(remainder(timer(),2),1)

Conclusion convert_to_tables

 Expand Expand on update
 Execute Run on clock tick

Directed Graphs

Example



Example

Dataset:

edge (a , b)

edge (b , d)

edge (b , e)

Operations:

copy (X , Y) - copy outgoing links from X to Y.

invert (X) - reverse the outgoing arcs from X.

insert (X , Y) - add arc from X to Y and all Y successors.

copy(X, Y)

Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e)\}$

Operation Definition:

$\text{copy}(X, Y) :: \text{edge}(X, Z) \implies \text{edge}(Y, Z)$

Action: $\text{copy}(b, c)$

Positive Updates: $\{\text{edge}(c, d), \text{edge}(c, e)\}$

Negative Updates: $\{\}$

Result:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(c, d), \text{edge}(c, e)\}$

invert(X)

Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(c, d), \text{edge}(c, e)\}$

Operation Definition:

$\text{invert}(X)$ - reverse the outgoing arcs from x .

invert(X)

Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(c, d), \text{edge}(c, e)\}$

Operation Definition:

$\text{invert}(X) :: \text{edge}(X, Y) \implies \sim\text{edge}(X, Y) \ \& \ \text{edge}(Y, X)$

Action: $\text{invert}(c)$

Positive Updates: $\{\text{edge}(d, c), \text{edge}(e, c)\}$

Negative Updates: $\{\text{edge}(c, d), \text{edge}(c, e)\}$

Result:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(d, c), \text{edge}(e, c)\}$

insert(X, Y)

Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(d, c), \text{edge}(e, c)\}$

Operation Definition:

$\text{insert}(X, Y)$ - add arc from X to Y and all successors of Y.

insert(X,Y)

Dataset:

$\{\text{edge}(a,b), \text{edge}(b,d), \text{edge}(b,e), \text{edge}(d,c), \text{edge}(e,c)\}$

Operation Definition:

$\text{insert}(X,Y) :: \text{edge}(X,Y)$

$\text{insert}(X,Y) :: \text{edge}(Y,Z) \implies \text{insert}(X,Z)$

Action: $\text{insert}(w,b)$

Expansion:

$\{\text{insert}(w,b), \text{edge}(w,b), \text{insert}(w,d), \text{insert}(w,e),$
 $\text{edge}(w,d), \text{edge}(w,e), \text{insert}(w,c), \text{edge}(w,c)\}$

Negative Updates: $\{\}$

Positive Updates:

$\{\text{edge}(w,b), \text{edge}(w,d), \text{edge}(w,e), \text{edge}(w,c)\}$

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
```

Execute

Action |

Expand Expand on update

Execute Run on clock tick

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
```

Execute

Action `copy(b,c)`

Expand Expand on update
Execute Run on clock tick

```
edge(c,d)
edge(c,e)
```

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
edge(c,d)
edge(c,e)
```

Execute

Action `copy(b,c)`

Expand Expand on update

Execute Run on clock tick

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
edge(c,d)
edge(c,e)
```

Execute

Action invert(c)

Expand Expand on update

Execute Run on clock tick

```
~edge(c,d)
edge(d,c)
~edge(c,e)
edge(e,c)
```

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
edge(d,c)
edge(e,c)
```

Execute

Action invert(c)

Expand Expand on update

Execute Run on clock tick

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
edge(d,c)
edge(e,c)
```

Execute

Action insert(w,b)

Expand Expand on update
Execute Run on clock tick

```
edge(w,b)
edge(w,d)
edge(w,c)
edge(w,e)
```

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

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File Dataset Channel Ruleset Operation Settings

Lambda

Save Revert Sort

```
edge(a,b)
edge(b,d)
edge(b,e)
edge(d,c)
edge(e,c)
edge(w,b)
edge(w,d)
edge(w,c)
edge(w,e)
```

Execute

Action insert(w,b)

Expand Expand on update

Execute Run on clock tick

Library

Save Revert

```
copy(X,Y) :: edge(X,Z) ==> edge(Y,Z)

invert(X) :: edge(X,Y) ==> ~edge(X,Y) & edge(Y,X)

insert(X,Y) :: edge(X,Y)
insert(X,Y) :: edge(Y,Z) ==> insert(X,Z)
```

Display a menu

View Materialization

Materialized Views

A **materialized view** is a view relation that is stored explicitly in the database.

Ruleset:

```
grandparent(X,Z) :- parent(X,Y) & parent(Y,Z)
```

Base Data:

```
parent(art,bob)  
parent(art,bea)  
parent(bob,cal)  
parent(bob,cam)  
parent(bea,cat)  
parent(bea,coe)
```

Materialized View:

```
grandparent(art,cal)  
grandparent(art,cam)  
grandparent(art,cat)  
grandparent(art,coe)
```

Analysis

Ruleset

```
s(X,Y,Z) :- r(X) & r(Y) & r(Z)
r(X) :- p(X,Y) & p(Y,Z)
```

Computation Cost for s:

r computed multiple times

for $n=2$, unifications = 1242

for $n=3$, unifications = 41636

where n is the number of objects

Storage for p:

$O(n^2)$ facts stored

Example with s Materialized

Ruleset

```
s(X, Y, Z) :- r(X) & r(Y) & r(Z)
r(X) :- p(X, Y) & p(Y, Z)
```

Computation Cost for s

for $n=2$, unifications = 8

for $n=3$, unifications = 27

Storage for s

$O(n^3)$ in worst case

Example with r Materialized

Ruleset

$$s(X, Y, Z) \text{ :- } r(X) \ \& \ r(Y) \ \& \ r(Z)$$
$$r(X) \text{ :- } p(X, Y) \ \& \ p(Y, Z)$$

Computation Cost for s

for $n=2$, unifications = 15

for $n=3$, unifications = 40

where n is the number of objects

Computation Cost for r

for $n=2$, unifications = 17

for $n=3$, unifications = 55

Storage for r

$O(n)$ in worst case

Materialization

Naive Approach:

On receipt of update,
discard materialized data and
rematerialize.

Differential Approach:

Write operation definitions for the materialized view.
Remove rules defining the materialized view.
On receipt of update, run newly defined operations.

Example

Old View Definitions

$s(X, Y, Z) :- r(X) \ \& \ r(Y) \ \& \ r(Z)$

$r(X) :- p(X, Y) \ \& \ p(Y, Z)$

Operation Definitions

$+p(X, Y) :: p(X, Y)$

$+p(X, Y) :: p(Y, Z) ==> r(X)$

$+p(Y, Z) :: p(X, Y) ==> r(X)$

Exercise for viewer: What are the deletion rules?

New View Definition (remove definition of r)

$s(X, Y, Z) :- r(X) \ \& \ r(Y) \ \& \ r(Z)$

Execution

Old Dataset: $\{p(a, b), p(b, a), r(a), r(b)\}$

Rules:

$s(X, Y, Z) :- r(X) \ \& \ r(Y) \ \& \ r(Z)$

$+p(X, Y) :: p(X, Y)$

$+p(X, Y) :: p(Y, Z) ==> r(X)$

$+p(Y, Z) :: p(X, Y) ==> r(X)$

Change Request: $+p(d, c)$

Update: $\{p(d, c)\}$

New Dataset:

$\{p(a, b), p(b, a), p(d, c), r(a), r(b)\}$

Execution

Old Dataset: $\{p(a, b), p(b, a), p(d, c), r(a), r(b)\}$

Rules:

$s(X, Y, Z) :- r(X) \ \& \ r(Y) \ \& \ r(Z)$

$+p(X, Y) :: p(X, Y)$

$+p(X, Y) :: p(Y, Z) ==> r(X)$

$+p(Y, Z) :: p(X, Y) ==> r(X)$

Change Request: $+p(c, d)$

Update: $\{p(c, d), r(c), r(d)\}$

New Dataset:

$\{p(a, b), p(b, a), p(c, d), p(d, c),$
 $r(a), r(b), r(c), r(d)\}$

Update Through Views

View Definition

$r(X) :- p(X)$

$r(X) :- q(X)$

Request: $+r(\text{bob})$

Positive Update:

$\{p(\text{bob})\}?$

$\{q(\text{bob})\}?$

$\{p(\text{bob}), q(\text{bob})\}?$ *What if p is dead and q is alive?*

$\{\}?$

Update Policies

View Definition

$r(X) \text{ :- } p(X)$

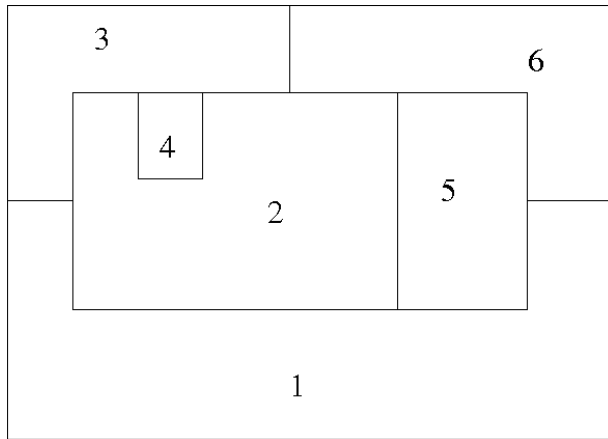
$r(X) \text{ :- } q(X)$

Update Rule

$+r(X) \text{ :: } \sim q(X) \implies p(X)$

Constraint Propagation

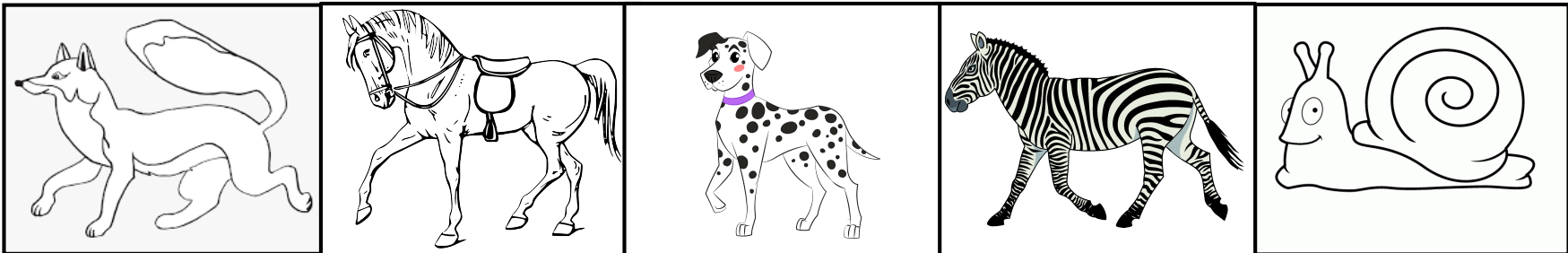
Constraint Satisfaction Problems



SEND
+MORE

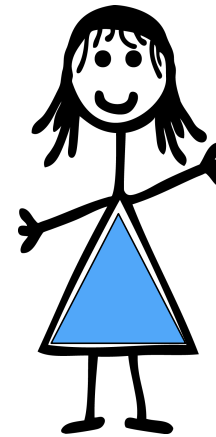
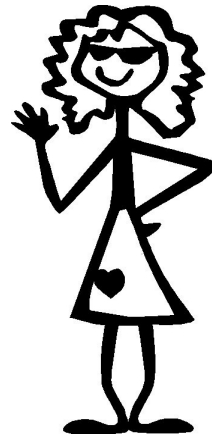
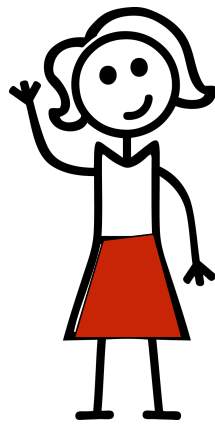
MONEY

	6		1	4		5	
		8	3		5	6	
2							1
8			4	7			6
		6				3	
7			9	1			4
5							2
		7	2		6	9	
	4		5		8		7



Problem Statement

Ruby Red, Willa White, Betty Blue are having lunch.
One is wearing a red skirt, one white, one blue.

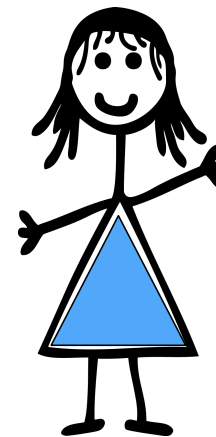
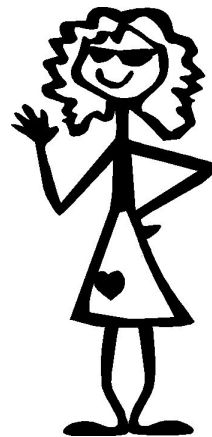
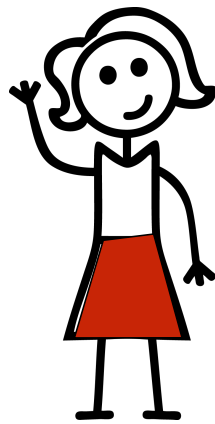


Betty to woman in white skirt: *Have you noticed that our skirts have colors different from our names?*

Which woman is wearing which skirt?

Problem Statement

Ruby Red, Willa White, Betty Blue are having lunch.
One is wearing a red skirt, one white, one blue.



Betty to woman in white skirt: *Have you noticed that our skirts have colors different from our names?*

Why is this problem so easy to solve?

Vocabulary

Symbols: red, white, blue, v, x

Unary predicates Predicates:

$\text{color}(C)$ is true if C is a color.

Ternary Predicates:

$c(P, C, v)$ is true if person P is wearing color C.

$c(P, C, x)$ is true if person P is not wearing color C.

Operations:

c1 - No person is wearing same color as name.

c2 - Betty Blue spoke to person in white.

c3 - Every person has one color and vice versa.

Colors

```
color(red)
```

```
color(white)
```

```
color(blue)
```

First Given

No person is wearing same color as name.

`c1 :: color(C) ==> c(C,C,x)`

Second Given

Betty Blue spoke to person in white.

```
c2 :: c(blue,white,x)
```


Existence Constraints

Every person is wearing some color.

`c3 ::`

```
c(P,C1,x) & c(P,C2,x) & color(C3) & mutex(C1,C2,C3)
==> c(P,C3,v)
```

Every color is worn.

`c3 ::`

```
c(P1,C,x) & c(P2,C,x) & color(P3) & mutex(P1,P2,P3)
==> c(P3,C,v)
```

Uniqueness Constraints

Nobody is wearing two colors.

`c3 ::`

```
c(P,C1,v) & color(C2) & distinct(C1,C2)  
==> c(P,C2,x)
```

No color is worn by two people.

`c3 ::`

```
c(P1,C,v) & color(P2) & distinct(P1,P2)  
==> c(P2,C,x)
```

Initial State

Color

	r	w	b
Person			
r			
w			
b			

First Given

No person is wearing same color as name.

`c1 :: color(C) ==> c(C,C,x)`

	r	w	b
r			
w			
b			



	r	w	b
r	X		
w		X	
b			X

Second Given

Betty Blue spoke to person in white.

```
c2 :: c(blue,white,x)
```

	r	w	b
r	X		
w		X	
b			X



	r	w	b
r	X		
w		X	
b		X	X

Existence Constraints

c3 ::

$c(P, C1, x) \ \& \ c(P, C2, x) \ \& \ color(C3) \ \& \ mutex(C1, C2, C3)$
 $\implies c(P, C3, v)$

c3 ::

$c(P1, C, x) \ \& \ c(P2, C, x) \ \& \ human(P3) \ \& \ mutex(P1, P2, P3)$
 $\implies c(P3, C, v)$

	r	w	b
r	X		
w		X	
b		X	X



	r	w	b
r	X	V	
w		X	
b	V	X	X

Uniqueness Constraints

c3 ::

$c(P, C1, v) \ \& \ color(C2) \ \& \ distinct(C1, C2)$
 $\implies c(P, C2, x)$

c3 ::

$c(P1, C, v) \ \& \ person(P2) \ \& \ distinct(P1, P2)$
 $\implies c(P2, C, x)$

	r	w	b
r	X	V	
w		X	
b	V	X	X



	r	w	b
r	X	V	X
w	X	X	
b	V	X	X

Existence Constraints

c3 ::

$c(P, C1, x) \ \& \ c(P, C2, x) \ \& \ color(C3) \ \& \ mutex(C1, C2, C3)$
 $\implies c(P, C3, v)$

c3 ::

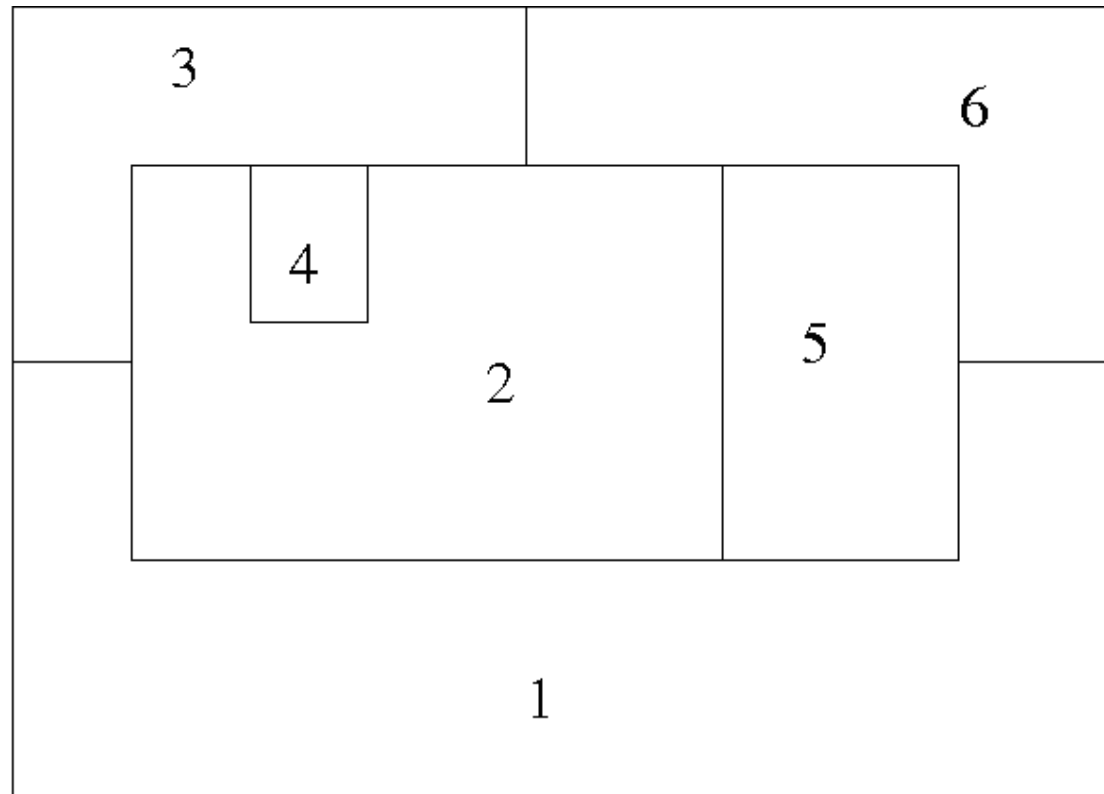
$c(P1, C, x) \ \& \ c(P2, C, x) \ \& \ human(P3) \ \& \ mutex(P1, P2, P3)$
 $\implies c(P3, C, v)$

	r	w	b
r	X	V	X
w	X	X	
b	V	X	X



	r	w	b
r	X	V	X
w	X	X	V
b	V	X	X

Map Coloring



Sudoku

	6		1		4		5	
		8	3		5	6		
2								1
8			4		7			6
		6				3		
7			9		1			4
5								2
		7	2		6	9		
	4		5		8		7	

